

The impact of post-quantum bugs

Daniel J. Bernstein

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- About 1/8 of “cryptographic” CVEs are severe.

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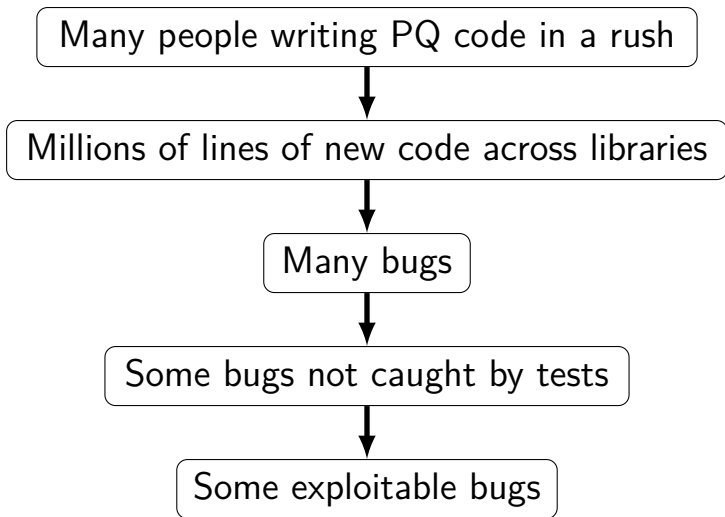
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Examples of bug announcements in PQ software:

2018 Dilithium	2019 Picnic	2024 SMAUG-T
2018 SPHINCS+	2019 qTESLA	2024 ML-DSA
2018 NTRU	2023 HAETAE	2026 ML-KEM
2019 Falcon	2023 PERK	2026 ML-DSA
2019 Picnic	2024 PALOMA	2026 ML-KEM
2019 NTRU	2024 NCC-Sign	2026 ML-DSA

Also many announcements of timing variations.

The PQ code ecosystem is getting worse



Case study: Dilithium (ML-DSA)

The first announced Dilithium bug

“The CRYSTALS Team” (Dilithium+Kyber), [2018](#):

“We are very grateful to Peter Pessl for notifying us of an implementation error in a randomness generator of our NIST submission. . . . the result of the bug was that the same randomness ended up being used for pairs of consecutive coefficients . . . This reuse of randomness can easily be exploited to recover the secret key and we thus emphasize that the software, in the state submitted to NIST, should not be used in any real application.”

One way to make this mistake yourself

Fixed `poly.c` obtains random `t0`, `t1`. Then:

```
if(t0 <= 2*GAMMA1 - 2)
    a[ctr++] = Q + GAMMA1 - 1 - t0;
if(t1 <= 2*GAMMA1 - 2 && ctr < len)
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(Original code actually used just one variable `t`.
Code intent: random `t`, use `t`, random `t`, use `t`.
Actual code: random `t`, random `t`, use `t`, use `t`.)

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Also checks that it can use the key it computed to forge signatures on attacker-chosen messages.

Details: <https://cr.yp.to/papers.html#mldsa>

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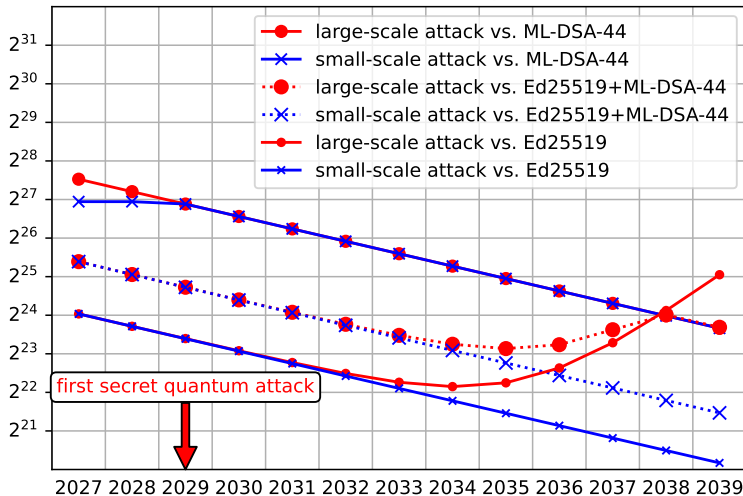
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Estimate 2^{30} target ML-DSA keys.

Account for limits on attacker CPU time.

See paper for more

Estimated number of broken keys in each year



The future

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(These also **catch** many timing leaks.)
6. Try using formal-verification tools. Examples:
CryptoLine, **s2n-bignum**, **saferewrite**, **sortverif**.
But **beware** overconfident claims of verification.