

A subfield-logarithm attack
against ideal lattices,
part 1: the number-field sieve

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Sieving small integers $i > 0$
using primes 2, 3, 5, 7:

1				
2	2			
3		3		
4	2 2			
5			5	
6	2	3		
7				7
8	2 2 2			
9		3 3		
10	2		5	
11				
12	2 2	3		
13				
14	2			7
15		3	5	
16	2 2 2 2			
17				
18	2	3 3		
19				
20	2 2		5	

etc.

Sieving i and $611 + i$ for small i
 using primes 2, 3, 5, 7:

1				
2	2			
3		3		
4	2 2			
5			5	
6	2	3		
7				7
8	2 2 2			
9		3 3		
10	2		5	
11				
12	2 2	3		
13				
14	2			7
15		3	5	
16	2 2 2 2			
17				
18	2	3 3		
19				
20	2 2		5	

612	2 2	3 3		
613				
614	2			
615		3	5	
616	2 2 2			7
617				
618	2	3		
619				
620	2 2		5	
621		3 3 3		
622	2			
623				7
624	2 2 2 2 3			
625			5 5 5 5	
626	2			
627		3		
628	2 2			
629				
630	2	3 3	5	7
631				

etc.

Have complete factorization of the “congruences” $i(611 + i)$ for some i 's.

$$14 \cdot 625 = 2^1 3^0 5^4 7^1.$$

$$64 \cdot 675 = 2^6 3^3 5^2 7^0.$$

$$75 \cdot 686 = 2^1 3^1 5^2 7^3.$$

$$14 \cdot 64 \cdot 75 \cdot 625 \cdot 675 \cdot 686 \\ = 2^8 3^4 5^8 7^4 = (2^4 3^2 5^4 7^2)^2.$$

$$\gcd\{611, 14 \cdot 64 \cdot 75 - 2^4 3^2 5^4 7^2\} \\ = 47.$$

$$611 = 47 \cdot 13.$$

Why did this find a factor of 611?

Was it just blind luck:

$$\gcd\{611, \text{random}\} = 47?$$

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No.

By construction 611 divides $s^2 - t^2$

where $s = 14 \cdot 64 \cdot 75$

and $t = 2^4 3^2 5^4 7^2$.

So each prime > 7 dividing 611 divides either $s - t$ or $s + t$.

Not terribly surprising

(but not guaranteed in advance!)

that one prime divided $s - t$

and the other divided $s + t$.

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 $(1, 0, 4, 1)$, $(6, 3, 2, 0)$, $(1, 1, 2, 3)$
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Yes. The exponent vectors $(1, 0, 4, 1)$, $(6, 3, 2, 0)$, $(1, 1, 2, 3)$ happened to have sum $0 \pmod 2$.

But we didn't need this luck!

Given long sequence of vectors, quickly find nonempty subsequence with sum $0 \pmod 2$.

This is linear algebra over \mathbf{F}_2 .

Guaranteed to find subsequence

if number of vectors

exceeds length of each vector.

e.g. for $n = 671$:

$$1(n + 1) = 2^5 3^1 5^0 7^1;$$

$$4(n + 4) = 2^2 3^3 5^2 7^0;$$

$$15(n + 15) = 2^1 3^1 5^1 7^3;$$

$$49(n + 49) = 2^4 3^2 5^1 7^2;$$

$$64(n + 64) = 2^6 3^1 5^1 7^2.$$

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\mathbf{F}_2 -kernel of exponent matrix is

gen by $(0\ 1\ 0\ 1\ 1)$ and $(1\ 0\ 1\ 1\ 0)$;

e.g., $1(n + 1)15(n + 15)49(n + 49)$

is a square.

Plausible conjecture: \mathbf{Q} sieve can separate the odd prime divisors of any n , not just 611.

Given n and parameter y :

1. Try to fully factor $i(n + i)$ into products of primes $\leq y$ for $i \in \{1, 2, 3, \dots, y^2\}$.

2. Look for nonempty set of i 's with $i(n + i)$ completely factored and with $\prod_i i(n + i)$ square.

3. Compute $\gcd\{n, s - t\}$ where $s = \prod_i i$ and $t = \sqrt{\prod_i i(n + i)}$.

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Let's aim for number of completely factored congruences to exceed length of each vector, guaranteeing a square.

(This is somewhat pessimistic; smaller numbers usually work.)

Vector length $\approx y/\log y$.

Will there be $> y/\log y$ completely factored congruences out of y^2 congruences?

What's chance of random $i(n+i)$ being y -**smooth**, i.e., completely factored into primes $\leq y$?

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Consider, e.g., $y = \lfloor n^{1/10} \rfloor$.

Uniform random integer in $[1, y^2]$ has y -smoothness chance ≈ 0.306 ;
uniform random integer in $[1, n]$ has chance $\approx 2.77 \cdot 10^{-11}$.

Plausible conjecture:

y -smoothness chance of $i(n+i)$ is $\approx 8.5 \cdot 10^{-12}$.

Find $\approx 8.5 \cdot 10^{-12} y^2$

fully factored congruences.

If $n \geq 2^{340}$ and $y = \lfloor n^{1/10} \rfloor$ then $8.5 \cdot 10^{-12} y^2 > 3y/\log y$, and approximations seem fairly close, so conjecturally the **Q** sieve will find a square.

Find many independent squares with negligible extra effort.

If gcd turns out to be 1, try the next square.

Conjecturally always works:
splits odd n into
prime-power factors.

How about $y \approx n^{1/u}$

for larger u ?

Uniform random integer in $[1, n]$

has $n^{1/u}$ -smoothness chance

roughly u^{-u} .

Plausible conjecture:

Q sieve succeeds

with $y = \lfloor n^{1/u} \rfloor$

for all $n \geq u^{(1+o(1))u^2}$;

here $o(1)$ is as $u \rightarrow \infty$.

How about

letting u grow with n ?

Given n , try sequence of y 's

in geometric progression

until **Q** sieve works;

e.g., increasing powers of 2.

Plausible conjecture: final $y \in$

$$\exp \sqrt{\left(\frac{1}{2} + o(1)\right) \log n \log \log n},$$

$$u \in \sqrt{(2 + o(1)) \log n / \log \log n}.$$

Cost of **Q** sieve is a power of y ,

hence subexponential in n .

More generally, if $y \in$
 $\exp \sqrt{\left(\frac{1}{2c} + o(1)\right) \log n \log \log n}$,
conjectured y -smoothness chance
is $1/y^{c+o(1)}$.

Find enough smooth congruences
by changing the range of i 's:
replace y^2 with $y^{c+1+o(1)} =$
 $\exp \sqrt{\left(\frac{(c+1)^2 + o(1)}{2c}\right) \log n \log \log n}$.

Increasing c past 1

increases number of i 's but
reduces linear-algebra cost.

So linear algebra never dominates
when y is chosen properly.

Improving smoothness chances

Smoothness chance of $i(n + i)$ degrades as i grows.

Smaller for $i \approx y^2$ than for $i \approx y$.

Crude analysis: $i(n + i)$ grows.

$\approx yn$ if $i \approx y$;

$\approx y^2n$ if $i \approx y^2$.

More careful analysis:

$n + i$ doesn't degrade, but

i is always smooth for $i \leq y$,

only 30% chance for $i \approx y^2$.

Can we select congruences to avoid this degradation?

Choose q , square of large prime.

Choose a “ q -sublattice” of i 's:

arithmetic progression of i 's

where q divides each $i(n + i)$.

e.g. progression $q - (n \bmod q)$,

$2q - (n \bmod q)$, $3q - (n \bmod q)$,

etc.

Check smoothness of

generalized congruence $i(n + i)/q$

for i 's in this sublattice.

e.g. check whether $i, (n + i)/q$ are

smooth for $i = q - (n \bmod q)$ etc.

Try many large q 's.

Rare for i 's to overlap.

e.g. $n = 314159265358979323$:

Original **Q** sieve:

i $n + i$

1 314159265358979324

2 314159265358979325

3 314159265358979326

Use 997^2 -sublattice,

$i \in 802458 + 994009\mathbf{Z}$:

i $(n + i)/997^2$

802458 316052737309

1796467 316052737310

2790476 316052737311

Crude analysis: Sublattices
eliminate the growth problem.
Have practically unlimited supply
of generalized congruences

$$(q - (n \bmod q)) \frac{n + q - (n \bmod q)}{q}$$

between 0 and n .

More careful analysis: Sublattices
are even better than that!

For $q \approx n^{1/2}$ have

$$i \approx (n + i)/q \approx n^{1/2} \approx y^{u/2}$$

so smoothness chance is roughly

$$(u/2)^{-u/2} (u/2)^{-u/2} = 2^u / u^u,$$

2^u times larger than before.

Even larger improvements
from changing polynomial $i(n+i)$.

“Quadratic sieve” (QS) uses
 $i^2 - n$ with $i \approx \sqrt{n}$;
have $i^2 - n \approx n^{1/2+o(1)}$,
much smaller than n .

“MPQS” improves $o(1)$
using sublattices: $(i^2 - n)/q$.
But still $\approx n^{1/2}$.

“Number-field sieve” (NFS)
achieves $n^{o(1)}$.

Generalizing beyond \mathbf{Q}

The \mathbf{Q} sieve is a special case of the number-field sieve.

Recall how the \mathbf{Q} sieve factors 611:

Form a square

as product of $i(i + 611j)$

for several pairs (i, j) :

$$14(625) \cdot 64(675) \cdot 75(686) \\ = 4410000^2.$$

$$\gcd\{611, 14 \cdot 64 \cdot 75 - 4410000\} \\ = 47.$$

The $\mathbf{Q}(\sqrt{14})$ sieve
factors 611 as follows:

Form a square

as product of $(i + 25j)(i + \sqrt{14}j)$

for several pairs (i, j) :

$$\begin{aligned} &(-11 + 3 \cdot 25)(-11 + 3\sqrt{14}) \\ &\quad \cdot (3 + 25)(3 + \sqrt{14}) \\ &= (112 - 16\sqrt{14})^2. \end{aligned}$$

Compute

$$s = (-11 + 3 \cdot 25) \cdot (3 + 25),$$

$$t = 112 - 16 \cdot 25,$$

$$\gcd\{611, s - t\} = 13.$$

Why does this work?

Answer: Have ring morphism
 $\mathbf{Z}[\sqrt{14}] \rightarrow \mathbf{Z}/611$, $\sqrt{14} \mapsto 25$,
since $25^2 = 14$ in $\mathbf{Z}/611$.

Apply ring morphism to square:

$$\begin{aligned} &(-11 + 3 \cdot 25)(-11 + 3 \cdot 25) \\ &\quad \cdot (3 + 25)(3 + 25) \\ &= (112 - 16 \cdot 25)^2 \text{ in } \mathbf{Z}/611. \end{aligned}$$

i.e. $s^2 = t^2$ in $\mathbf{Z}/611$.

Unsurprising to find factor.

Diagram of ring morphisms:

$$\begin{array}{ccc}
 \mathbf{Q}[x] & \xrightarrow{x \mapsto \sqrt{14}} & \mathbf{Q}[\sqrt{14}] = \mathbf{Q}(\sqrt{14}) \\
 \uparrow & & \uparrow \\
 \mathbf{Z}[x] & \xrightarrow{x \mapsto \sqrt{14}} & \mathbf{Z}[\sqrt{14}] \\
 & & \downarrow \sqrt{14} \mapsto 25 \\
 & & \mathbf{Z}/611
 \end{array}$$

$\mathbf{Z}[x]$ uses poly arithmetic on
 $\{i_0x^0 + i_1x^1 + \dots : \text{all } i_m \in \mathbf{Z}\}$;
 $\mathbf{Z}[\sqrt{14}]$ uses \mathbf{R} arithmetic on
 $\{i_0 + i_1\sqrt{14} : i_0, i_1 \in \mathbf{Z}\}$;
 $\mathbf{Z}/611$ uses arithmetic mod 611
 on $\{0, 1, \dots, 610\}$.

Generalize from $(x^2 - 14, 25)$
to (f, m) with irred $f \in \mathbf{Z}[x]$,
 $m \in \mathbf{Z}$, $f(m) \in n\mathbf{Z}$.

Write $d = \deg f$,

$$f = f_d x^d + \cdots + f_1 x^1 + f_0 x^0.$$

Can take $f_d = 1$ for simplicity,
but larger f_d allows
better parameter selection.

Pick $\alpha \in \mathbf{C}$, root of f .

Then $f_d \alpha$ is a root of
monic $g = f_d^{d-1} f(x/f_d) \in \mathbf{Z}[x]$.

$$\mathbf{Q}(\alpha) = \left\{ \begin{array}{l} r_0 + r_1\alpha + r_2\alpha^2 + \\ \cdots + r_{d-1}\alpha^{d-1}. \\ r_0, \dots, r_{d-1} \in \mathbf{Q} \end{array} \right\}$$



$$\mathcal{O} = \left\{ \begin{array}{l} \text{algebraic integers} \\ \text{in } \mathbf{Q}(\alpha) \end{array} \right\}$$



$$\mathbf{Z}[f_d\alpha] = \left\{ \begin{array}{l} i_0 + i_1 f_d\alpha + \\ \cdots + i_{d-1} f_d^{d-1} \alpha^{d-1}. \\ i_0, \dots, i_{d-1} \in \mathbf{Z} \end{array} \right\}$$



$$f_d\alpha \mapsto f_d m$$

$$\mathbf{Z}/n = \{0, 1, \dots, n - 1\}$$

Build square in $\mathbf{Q}(\alpha)$ from
congruences $(i - jm)(i - j\alpha)$
with $i\mathbf{Z} + j\mathbf{Z} = \mathbf{Z}$ and $j > 0$.

Could replace $i - jx$ by
higher-deg irred in $\mathbf{Z}[x]$;
quadratics seem fairly small
for some number fields.

But let's not bother.

Say we have a square

$\prod_{(i,j) \in S} (i - jm)(i - j\alpha)$
in $\mathbf{Q}(\alpha)$; now what?

$$\prod (i - jm)(i - j\alpha) f_d^2$$

is a square in \mathcal{O} ,

ring of integers of $\mathbf{Q}(\alpha)$.

Multiply by $g'(f_d\alpha)^2$,

putting square root into $\mathbf{Z}[f_d\alpha]$:

compute r with $r^2 = g'(f_d\alpha)^2$.

$$\prod (i - jm)(i - j\alpha) f_d^2.$$

Then apply the ring morphism

$\varphi : \mathbf{Z}[f_d\alpha] \rightarrow \mathbf{Z}/n$ taking

$f_d\alpha$ to f_dm . Compute $\gcd\{n,$

$\varphi(r) - g'(f_dm) \prod (i - jm) f_d\}$.

In \mathbf{Z}/n have $\varphi(r)^2 =$

$$g'(f_dm)^2 \prod (i - jm)^2 f_d^2.$$

How to find square product
of congruences $(i - jm)(i - j\alpha)$?

Start with congruences for,
e.g., y^2 pairs (i, j) .

Look for y -smooth congruences:

y -smooth $i - jm$ and

y -smooth $f_d \text{norm}(i - j\alpha) =$

$$f_d i^d + \cdots + f_0 j^d = j^d f(i/j).$$

Find enough smooth congruences.

Perform linear algebra on

exponent vectors mod 2.

Exponent vectors have
many “rational” components,
many “algebraic” components,
a few “character” components.

One rational component
for each prime $p \leq y$.

Value $\text{ord}_p(i - jm)$.

One rational component for -1 .

Value 0 if $i - jm > 0$,

value 1 if $i - jm < 0$.

If $\prod(i - jm)$ is a square

then vectors add to 0

in rational components.

One algebraic component
for each pair (p, r) such that
 p is a prime $\leq y$;

$$f_d \notin p\mathbf{Z}; \text{ disc } f \notin p\mathbf{Z};$$

$$r \in \mathbf{F}_p; f(r) = 0 \text{ in } \mathbf{F}_p.$$

Value 0 if $i - jr \notin p\mathbf{Z}$;

otherwise $\text{ord}_p(j^d f(i/j))$.

This is the same as

the valuation of $i - j\alpha$

at the prime $p\mathcal{O} + (f_d\alpha - f_dr)\mathcal{O}$.

Recall that $i\mathbf{Z} + j\mathbf{Z} = \mathbf{Z}$,

so no higher-degree primes.

One character component
for each pair (p, r) with
 p in a short range above y .

Value 0 if $i - jr$ is a
square in \mathbf{F}_p , else 1.

If $\prod(i - j\alpha)$ is a square
then vectors add to 0
in algebraic components
and character components.

Conversely, consider vectors
adding to 0 in all components.

$\prod(i - jm)$ must be a square.

Is $\prod(i - j\alpha)$ a square?

Ideal $\prod(i - j\alpha)\mathcal{O}$ must be
square outside f_d disc f .

What about primes in f_d disc f ?

Even if ideal is square,

is square root principal?

Even if ideal is generated

by square of element,

does square equal $\prod(i - j\alpha)$?

Obstruction group is small,
conjecturally very small.

“(f_d disc f)-Selmer group.”

A few characters
suffice to generate dual,
forcing $\prod (i - j\alpha)$
to be a square.

Can be quite sloppy here;
easy to redo linear algebra
with more characters if
non-square is encountered.

Sublattices

Consider a sublattice of pairs (i, j) where q divides $j^d f(i/j)$.

Assume squarish lattice.

$(i - jm)j^d f(i/j)$
expands by factor $q^{(d+1)/2}$
before division by q .

Number of sublattice elements within any particular bound

on $(i - jm)j^d f(i/j)$
is proportional to $q^{-(d-1)/(d+1)}$.

Compared to just using $q = 1$,
conjecturally obtain $y^{4/(d+1)+o(1)}$
times as many congruences
by using sublattices for
all y -smooth integers $q \leq y^2$.

Separately consider
 $i - jm$ and $j^d f(i/j)/q$
for more precise analysis.

Limit congruences accordingly,
increasing smoothness chances.

Multiple number fields

Assume that $f + x - m \in \mathbf{Z}[x]$
is also irred.

Pick $\beta \in \mathbf{C}$, root of $f + x - m$.

Two congruences for (i, j) :

$$(i - jm)(i - j\alpha); (i - jm)(i - j\beta).$$

Expand exponent vectors to
handle both $\mathbf{Q}(\alpha)$ and $\mathbf{Q}(\beta)$.

Merge smoothness tests

by testing $i - jm$ first,

aborting if $i - jm$ not smooth.

Can use many number fields:

$$f + 2(x - m) \text{ etc.}$$