

The algorithm of Hastad, Vallée,  
Girault, Toffin, Coppersmith,  
Guruswami, Goldreich, Ron,  
Sudan, Durfee, Howgrave-Graham,  
and Boneh

D. J. Bernstein

University of Illinois at Chicago

## Input to the algorithm

$d \in \mathbb{Z}, d \geq 1;$

$f \in \mathbf{Q}[x], \deg f = d, f_d > 0;$

$H \in \mathbb{Z}, H \geq 1;$

$k \in \mathbb{Z}, k \geq 1;$

$m \in \mathbb{Z}, m \geq dk.$

Simple case:  $k = 1, m = 2d.$

## Goal of the algorithm

Would like to find all  $r \in \mathbb{Q}$   
such that  $f(r)$  has small height.

Algorithm finds all  $r \in \mathbb{Z}$   
such that  $|r| \leq H$   
and  $\gcd\{1, f(r)\}/f_d > G$ .

Here  $\log \log G \approx$   
 $\frac{1}{2}(\log \log(1/f_d) + \log \log((2H)^d))$ .

## The algorithm

Define  $L \subset \mathbf{Q}[x]$  as

$$\begin{aligned} & \mathbf{Z} + \mathbf{Z}x + \mathbf{Z}x^2 + \dots + \mathbf{Z}x^{d-1} \\ & + \mathbf{Z}f + \dots + \mathbf{Z}x^{d-1}f \\ & + \mathbf{Z}f^2 + \dots + \mathbf{Z}x^{d-1}f^2 \\ & + \dots \\ & + \mathbf{Z}f^{k-1} + \dots + \mathbf{Z}x^{d-1}f^{k-1} \\ & + \mathbf{Z}f^k + \dots + \mathbf{Z}x^{m-dk-1}f^k. \end{aligned}$$

Define  $|\varphi| = \sqrt{\sum_i (H^i \varphi_i)^2}$ .

$L$  is a rank- $m$  lattice  
under metric  $\varphi \mapsto |\varphi|$ .

$$\det L = H^{\frac{1}{2}m(m-1)} f_d^{km - \frac{1}{2}dk(k+1)}.$$

Use LLL to find nonzero  $\varphi \in L$   
with  $|\varphi| \leq 2^{(m-1)/2} (\det L)^{1/m}$ .

Print all rational roots of  $\varphi$ .

# Speed of the algorithm

Tolerable.

$\varphi(r) = 0$  if  $r \in \mathbf{Z}$ ,  $|r| \leq H$ , and  
 $\gcd\{1, f(r)\}/f_d > G$ .

Here  $G = m^{1/2k} (2H)^{(m-1)/2k} f_d^{-d(k+1)/2m}$ .

Proof:  $|\varphi(r)| \leq m^{1/2} |\varphi|$   
 $\leq m^{1/2} 2^{(m-1)/2} (\det L)^{1/m}$   
 $= G^k f_d^k < \gcd\{1, f(r)\}^k$ ;  
 but  $\varphi(r) \in \gcd\{1, f(r)\}^k \mathbf{Z}$ .

## Good choice of $m$

Assume  $1/f_d \geq (2H)^d$ .

Take  $m = \lceil \alpha d(k+1) \rceil$  where

$$\alpha = \sqrt{\log(1/f_d) / \log((2H)^d)}.$$

Then  $G \leq m^{1/2k} (2H)^{\alpha d(1+1/2k)}$ .

So  $\varphi(r) = 0$  if  $r \in \mathbb{Z}$ ,  $|r| \leq H$ ,

and  $\gcd\{1, f(r)\}/f_d > m^{1/2k} (2H)^{\alpha d(1+1/2k)}$ .

## Application: roots mod $n$

$$f = \frac{(x+31415926000)^3 - 35083765367852945}{38785285061353277}$$

$$d = 3, H = 500, k = 1, m = 5.$$

$$L = \mathbf{Z} + \mathbf{Z}x + \mathbf{Z}x^2 + \mathbf{Z}f + \mathbf{Z}xf.$$

$$G < 20076370177628953 < 1/f_d.$$

Find  $\varphi \in L$  with  $38785285061353277\varphi =$   
 $250x^4 + 4480597x^3 - 3789099173x^2$   
 $+ 1135485860787x - 172139635662493.$

Integer roots of  $\varphi$ : 467.

$$f(467) = \frac{31006276476301184532318266236618}{38785285061353277}$$
$$= 799434023167634.$$

Given  $n \geq 1$ ,  $H \geq 1$ ,  $d \geq 1$ ,  
 $p \in \mathbf{Z}[x]$ ,  $\deg p = d$ ,  $p_d = 1$ :

Can apply algorithm  
with  $f = p/n$ ,  $k = 1$ ,  $m = d + 1$   
to find all  $r \in \mathbf{Z}$  with  
 $|r| \leq H$  and  $p(r) \in n\mathbf{Z}$ ,  
if  $H < n^{2/d(d+1)} / 2(d + 1)^{1/d}$ .

(Hastad 1985; complicated dual:  
Vallée, Girault, Toffin 1988)

Can apply algorithm  
with  $f = p/n$ ,  $k \geq 1$ ,  $m = dk + d$   
to find all  $r \in \mathbb{Z}$  with  
 $|r| \leq H$  and  $p(r) \in n\mathbb{Z}$ ,  
if  $H < n^{k/(m-1)} / 2m^{1/(m-1)}$ .

(dual: Coppersmith 1996;  
Howgrave-Graham 1997)

## Application: high-power factors

$$f = \frac{(4349000 + x)^2}{1038397528952788140203}$$

$$d = 2, H = 500, k = 2, m = 9.$$

$$\begin{aligned} L = & \mathbf{Z} + \mathbf{Z}_x + \mathbf{Z}_f + \mathbf{Z}_{xf} + \mathbf{Z}_{f^2} \\ & + \mathbf{Z}_{xf^2} + \mathbf{Z}_{x^2f^2} + \mathbf{Z}_{x^3f^2} + \mathbf{Z}_{x^4f^2}. \end{aligned}$$

$$G < 4348500^2.$$

Find small  $\varphi \in L$ .

Integer roots of  $\varphi$ : 353.

$f(353) = 1/54892667$ .

Try to factor integers this way.

Sometimes faster than ECM.

(Boneh, Durfee,  
Howgrave-Graham 1999)

## Application: CRT with errors

$$f = \frac{x - 1800140090020646934}{9156001667401012567}$$

$$d = 1, H = 5000, k = 1, m = 3.$$

$$L = \mathbb{Z} + \mathbb{Z}f + \mathbb{Z}xf.$$

$$f(3277) = -8675309/44124979.$$

$9156001667401012567 =$

$11 \cdot 13 \cdot 17 \cdot \dots \cdot 59.$

$3277 \bmod 11, 13, 17, \dots, 59:$

$10, 1, 13, 9, 11, 0, 22,$

$21, 38, 9, 34, 44, 32.$

$1800140090020646934 \bmod \dots:$

$10, 1, 4, 9, 11, 28, 22,$

$4, 14, 9, 34, 44, 29.$

Given  $H \geq 1$ ,  $n \geq 2H$ ,  $u \in \mathbf{Z}$ :

Write  $\alpha = \sqrt{\log(n)/\log(2H)}$ .

Can apply algorithm with

$$f = (x - u)/n, k \geq 1,$$

$m = \lceil \alpha(k + 1) \rceil$  to find all  $r \in \mathbf{Z}$

with  $|r| \leq H$  and  $\gcd\{n, r - u\} > m^{1/2k}(2H)^{\alpha(1+1/2k)}$ .

(Boneh 2000;  
dual with slightly worse result:  
Goldreich, Ron, Sudan 1998)

Similarly for function fields.

(Guruswami, Sudan 1999;  
slightly worse result: Sudan 1997)

## Application: smoothness

$$50!f = x^2 + 6563806563806000x + 14289695657685151430824671$$

$$d = 2, H = 1000, k = 2, m = 15.$$

$$50!f(823) = 2^{15} 3^{10} 5^4 11^3 13^1 19^2 29^1 37^1 41^1 43^1.$$

(Boneh 2000)